

Dictionary-based model reduction for state estimation

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Context

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- Parametric equation in a Hilbert space U ,

$$\mathcal{F}(u, \xi) = 0,$$

$u \in U$ the **state to recover** and $\xi \in \mathcal{P} \subset \mathbb{R}^d$ an **unknown** parameter.

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of m continuous linear measurements

$\ell_1(u), \dots, \ell_m(u)$, i.e. $w = P_W u$

with $W := \text{span}\{R_U \ell_i\}_{i=1}^m$

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model reduction



Compute a recovery $A(w) \simeq u$



One space problem

Linear MOR

One space: formulation

- Approximate \mathcal{M} by a linear subspace V ,

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- Parameterized-Background Data-Weak (PBDW) [1]:

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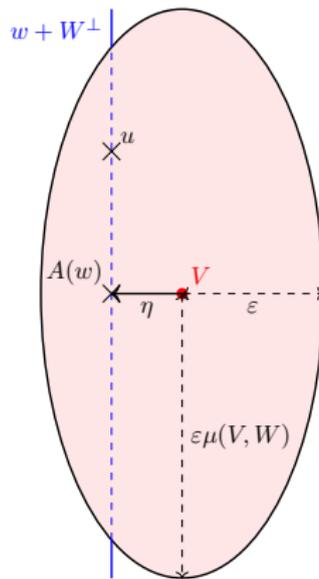
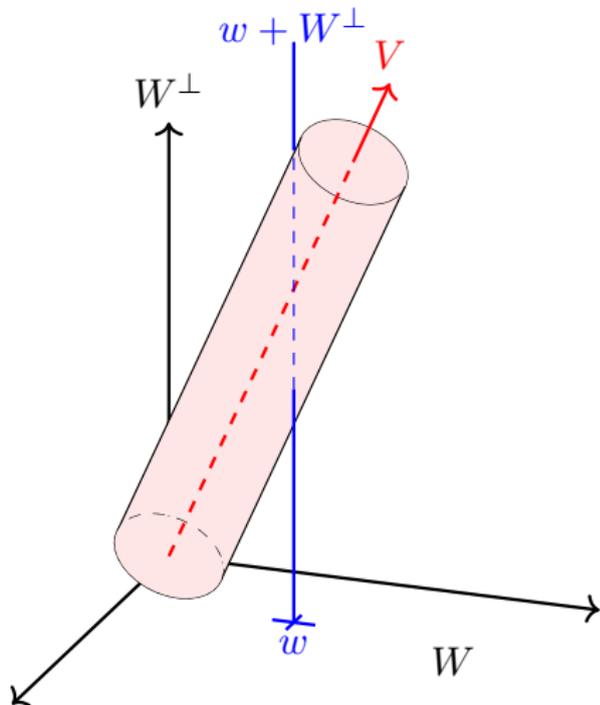
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- Sharp error bound [2]: $\|u - A_V(w)\| \leq \mu(V, W)\varepsilon$

$$\mu(V, W) := \beta(V, W)^{-1}, \quad \beta(V, W) := \inf_{v \in V} \sup_{w \in W} \frac{\langle v, w \rangle}{\|v\| \|w\|}$$

One space: Geometry



One space: Pros and cons

Pros

- Online efficiency $\mathcal{O}(n^3)$.
- No need to know ε .
- Optimal (worst case sens) when \mathcal{M} is a cylinder centered in V [2].

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Cons

- Requires $n \leq m$.
- Trade-off between ε and $\mu(V, W)$.
- Limited by the Kolmogorov m -width of \mathcal{M} ,

$$\varepsilon \geq d_m(\mathcal{M}).$$

Multi-space problem

Library-based MOR

Multi-space: library-based MOR

- Consider a library of spaces $\mathcal{L}_n^N := \{V_1, \dots, V_N\}$

$$\text{dist}\left(\bigcup_{k=1}^N V_k, \mathcal{M}\right) \leq \varepsilon, \quad \dim(V_k) \leq n \leq m.$$

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- New benchmark: non-linear Kolmogorov (n, N) -width [3]

$$d_n(\mathcal{M}, N) := \inf_{\#\mathcal{L}_n^N=N} \sup_{u \in \mathcal{M}} \min_{V \in \mathcal{L}_n^N} \min_{v \in V} \|u - v\|,$$

which is expected to decay much faster than

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- Aim for low ε with a low n , thus better stability.

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Each space $V \in \mathcal{L}_n^N$ gives a one-space estimate $A_V(w)$...

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How to select V^*
among this library ?

Multi-space: Selection

Idea [4]: select the "closest" to \mathcal{M}

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- Assume we have \mathcal{S} such that for any $v \in U$,

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- Select V^* as

$$V^* \in \operatorname{argmin}_{V \in \mathcal{L}_n^N} \mathcal{S}(A_V(w), \mathcal{M}).$$

Multi-space: Selection

Proposition (Near optimal selection [4])

Assuming that P_W is injective on \mathcal{M} and that $\mu(\mathcal{M}, W) < \infty$,

$$\|u - A_{k^*}(w)\| \leq 2 \frac{C}{c} \mu(\mathcal{M}, W) \min_{1 \leq k \leq N} \|u - A_k(w)\|,$$

$\mu(\mathcal{M}, W)$ reflects how well \mathcal{M} and W are aligned.

Dictionary-based multi-space

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Dictionary multi-space: Library

- Dictionary of K vectors (or snapshots),

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- Take as library $\mathcal{L}_m^N = \mathcal{L}_m(\mathcal{D}_K)$,

$$\mathcal{L}_m(\mathcal{D}_K) := \left\{ \sum_{k=1}^K x_k v^{(k)} : (x_k)_{k=1}^K \in \mathbb{R}^K, \|x\|_0 \leq m \right\},$$

Dictionary multi-space: Library

$\mathcal{L}_m(\mathcal{D}_K)$ is **large** \rightarrow low ε but not fully explorable.



Dictionary multi-space: Selection

- Consider the LASSO_α problem

$$\mathbf{x}^*(\alpha) \in \operatorname{argmin}_{\mathbf{x} \in \mathbb{R}^K} \frac{1}{2} \|\mathbf{C}\mathbf{x} - \mathbf{w}\|_2^2 + \alpha \|\mathbf{x}\|_1,$$

with $\alpha > 0$, $\mathbf{C} := (\langle w^{(i)}, v^{(j)} \rangle) \in \mathbb{R}^{m \times K}$ and $\mathbf{w} := (\langle w^{(i)}, u \rangle) \in \mathbb{R}^m$.

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- Use LARS_{α_0} to solve LASSO_{α_0} and generate the recoveries

$$(V_\alpha)_{\alpha \geq \alpha_0}, \quad V_\alpha := \operatorname{span}\{v^{(i)} : \mathbf{x}^*(\alpha)_i \neq 0\}.$$

with the associated recoveries $A_\alpha(w) := A_{V_\alpha}(w)$

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with the associated recoveries $A_\alpha(w) := A_{V_\alpha}(w)$

- Use \mathcal{S} to select $\alpha_{\mathcal{S}}$ among this **smaller** family.

Dictionary multi-space: Selection

Proposition (Near optimal selection [4])

Assuming that P_W is injective on \mathcal{M} and that $\mu(\mathcal{M}, W) < \infty$,

$$\|u - A_{\alpha_S}(w)\| \leq 2 \frac{C}{c} \mu(\mathcal{M}, W) \min_{\alpha > \alpha_0} \|u - A_\alpha(w)\|,$$

Parameterized PDEs

Offline-online decomposition for dictionary-based multispace

Parameterized PDEs: Framework

- Assume affinely parameterized operator equation (e.g. PDEs)

$$B(\xi)u = f \quad \text{with} \quad B(\xi) = B^{(0)} + \sum_{q=1}^{m_B} \xi_q B^{(q)}$$

with $B^{(q)} : U \rightarrow U'$ linear operators.

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- Assume **uniform control** on the singular values of $B(\xi)$

$$0 < c \leq \min_{v \in U} \frac{\|B(\xi)v\|}{\|v\|} \leq \max_{v \in U} \frac{\|B(\xi)v\|}{\|v\|} \leq C < \infty,$$

Parameterized PDEs: Surrogate distance

- Consider the residual-based distance

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$$\mathcal{S}(v, \mathcal{M}) := \min_{\xi \in \mathcal{P}} \|B(\xi)v - f\|, \quad v \in U$$

- $\mathcal{S}(v, \mathcal{M})$ is **computable** by solving a l.s. system with d dofs,

$$\mathcal{S}(v, \mathcal{M}) := \min_{\xi \in \mathcal{P}} \|G(v)\xi - g(v)\|$$

with

$$G(v) := \left(B^{(1)}v \mid \dots \mid B^{(d)}v \right) \quad \text{and} \quad g(v) := f - B^{(0)}v.$$

Parameterized PDEs: Offline-online decomposition

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- **Offline:** Heavy pre-computations independently on w .
- **Online:** Fast computation of $A_{\alpha_S}(w)$.

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- Instead, we use a random matrix $\Theta \in \mathbb{R}^{k \times \mathcal{N}}$ to compute

$$\mathcal{S}^\Theta(v, \mathcal{M}) := \min_{\xi \in \mathcal{P}} \|\Theta(B(\xi)v - f)\| = \min_{\xi \in \mathcal{P}} \|G^\Theta(v)\xi - g^\Theta(v)\|$$

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Proposition (Derived from [5])

With $k = \mathcal{O}\left(\epsilon^{-2} (d + \log(\delta^{-1}))\right)$, for any $v \in U$, with probability at least $1 - \delta$ we have

$$\sqrt{1 - \epsilon} \mathcal{S}(v, \mathcal{P}) \leq \mathcal{S}^\Theta(v, \mathcal{P}) \leq \sqrt{1 + \epsilon} \mathcal{S}(v, \mathcal{P}).$$

Parameterized PDEs: Offline

- Write $A_\alpha(w) = \mathbf{U}\mathbf{a}$ with $\mathbf{a} = \mathbf{a}(w) \in \mathbb{R}^{m+K}$ and $\mathbf{U} = (\mathbf{W} \mid \mathbf{V})$.

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- Write the l.s. terms to compute \mathcal{S}^Θ as

$$G^\Theta(A_\alpha(w)) := \left(\Theta B^{(1)} \mathbf{U}\mathbf{a} \mid \dots \mid \Theta B^{(d)} \mathbf{U}\mathbf{a} \right)$$
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- **Offline:** compute $\underbrace{\mathbf{C}}_{\mathcal{O}(mKN)} \in \mathbb{R}^{m \times K}$ and $\underbrace{\Theta B^{(q)} \mathbf{U}}_{\mathcal{O}(KN \log(\mathcal{N}))} \in \mathbb{R}^{k \times (m+K)}$

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- Total offline cost (without snapshot computation)

$$\mathcal{O}(mKN) + dKN \log(\mathcal{N})$$

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- **Step 2:** Prepare the l.s. systems, each costing $\mathcal{O}(kmd)$.
- **Step 3:** Solve them, each costing $\mathcal{O}(kd^2)$.
- Total online cost with $k = \mathcal{O}(d)$

$$\mathcal{O}\left(\underbrace{m^2 K}_{\text{LARS}} + \underbrace{md^2 K}_{\text{prepare l.s.}} + \underbrace{d^3 K}_{\text{solve l.s.}}\right).$$

Numerical Example

Finite Element space $U = \mathcal{H}_{\Gamma_D}^1(\Omega)$ embedded with $\|\nabla \cdot\|_{\mathcal{L}^2(\Omega)}$

Numerical: Thermal block

$$\mathcal{N} \sim 8\,000 \quad \text{and} \quad \begin{cases} -\nabla \cdot (\kappa \nabla u) = 1 & \text{in } \Omega, \\ u = 0 & \text{on } \Gamma_D, \\ \kappa = \xi_i & \text{in } \Omega_i, \quad 1 \leq i \leq 9, \end{cases}$$

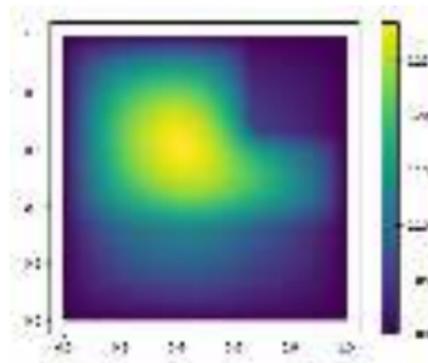
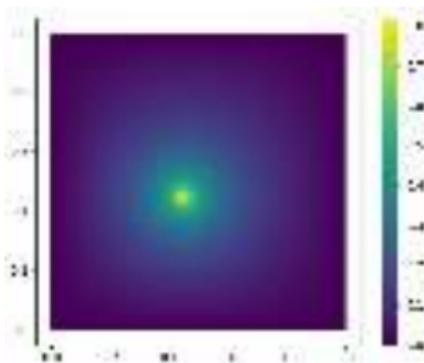
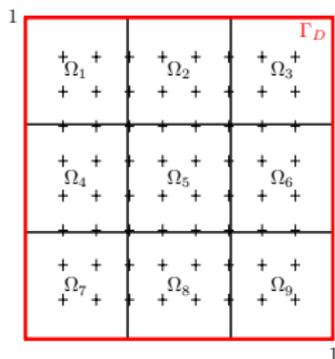


Figure: Left: geometry, with sensors locations (crosses). Middle: Riesz representer of a sensor. Right: a snapshot.

Numerical: Thermal block

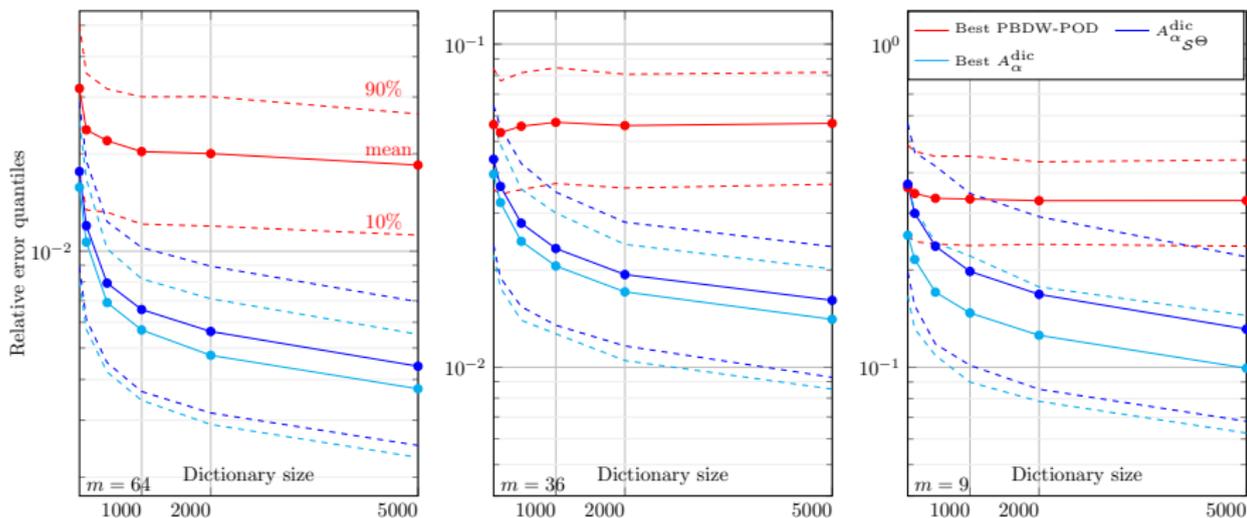
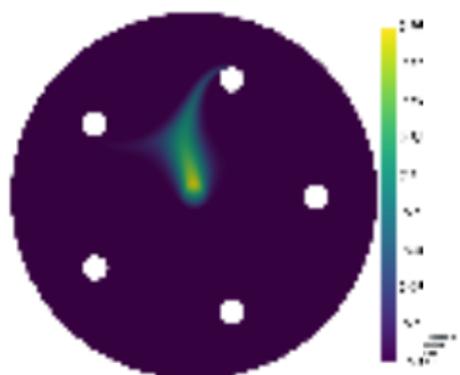
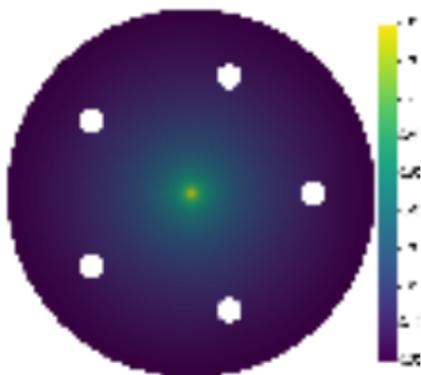
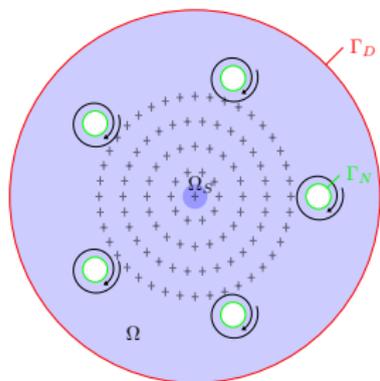


Figure: Evolution of the recovery errors in U -norm, on 500 test snapshots, with growing dictionary sizes.

Numerical: Advection diffusion

$$\mathcal{N} \sim 150\,000 \quad \text{and} \quad \begin{cases} -0.01\Delta u + \mathcal{V}(\xi) \cdot \nabla u = \frac{100}{\pi} \mathbb{1}_{\Omega_S} & \text{in } \Omega, \\ u = 0 & \text{on } \Gamma_D, \\ n \cdot \nabla u = 0 & \text{on } \Gamma_N, \end{cases}$$

$$\mathcal{V}(\xi) = \sum_{i=1}^5 \frac{1}{\|x - x^{(i)}\|} \left(\xi_i e_r(x^{(i)}) + \xi_{i+5} e_\theta(x^{(i)}) \right)$$



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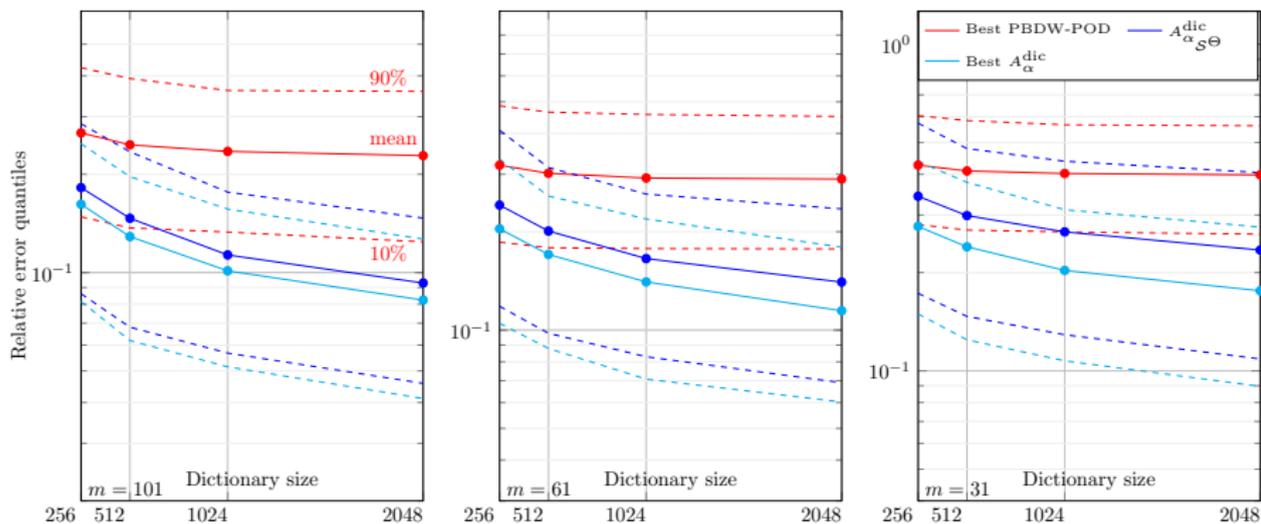


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Conclusion and perspectives

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Perspectives

- Noisy framework ?
- Bi-dictionary approach ?

Thank you !

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